

## Chapter 9. Planarity

**Definitions.** A *plane graph*  $G$  or a *planar embedding* of  $G$  is a drawing of  $G$  on the plane in such a way that no two edges meet, except at their common ends. Graphs that do admit such an embedding are called *planar*; ones that don't are called *nonplanar*.

A *girth* of a graph is the length of any smallest cycle if any. An acyclic graph has girth  $\infty$ . Thus,  $3 \leq \text{girth } G \leq \infty$  for any (simple) graph  $G$ .

Examples of planar graphs are paths, cycles, trees, and the complete bipartite graphs  $K_{2,k}$ .

Consider a plane graph. A *region* is a maximal connected area that remains when the edges & vertices are removed from the plane. The *boundary* of a region is the vertices and edges touching the region.

**Theorem** (Jordan Curve Theorem). *A simple closed curve partitions the plane into two regions: a bounded interior region and an unbounded exterior region.*

**Lemma** (Lemma A). *Any bridge is the boundary of exactly one region. Deleting a bridge (and any resulting isolated vertex) from a plane graph does not change the number of regions. Any nonbridge edge is the boundary of exactly two regions. Deleting a nonbridge edge from a plane graph decreases the number of regions by one.*

**Theorem** (Euler Identity, Theorem 9.1 of CZ). *If  $G$  is a connected plane graph of order  $n$ , size  $m$ , and  $r$  regions, then  $n - m + r = 2$ .*

*Proof.* We prove by induction on the number of cycles in  $G$ . If  $G$  has 0 cycle, then  $G$  is a tree since  $G$  is connected by assumption. Thus,  $m = n - 1$  and  $r = 1$ . Therefore,  $n - m + r = n - (n - 1) + 1 = 2$  and the result holds in the base case.

Now let  $G$  have  $k$  cycles, where  $k > 0$ , and assume inductively that any connected plane graph having fewer than  $k$  cycles satisfies the statement of the theorem. Let  $e$  be an edge belonging to some cycle of  $G$ . The plane graph  $G - e$  has  $n$  vertices,  $m - 1$  edges, and  $r - 1$  regions. Moreover,  $G - e$  is connected and has fewer than  $k$  cycles. Therefore, by the inductive hypothesis the result holds for  $G - e$ , i.e.,  $n - (m - 1) + (r - 1) = 2$ . This implies that  $n - m + r = 2$ , so the result holds for  $G$  as well.  $\square$

**Theorem** (Generalization of Theorem 9.2 of CZ). *Let  $g$  be a fixed integer  $\geq 3$ . If  $G$  is a planar graph of order  $n$ , size  $m$ , girth  $\geq g$ , and  $n \geq (g+2)/2$ , then  $m \leq \frac{g(n-2)}{g-2}$ .*

*Proof.* Note that for any planar graph  $G_1$ , there is a connected, planar graph  $G_2$  that is a supergraph of  $G_1$ . Thus we may assume that  $G$  is connected.

First, assume  $G$  has  $< g$  edges. Then  $G$  is acyclic since it has girth  $\geq g$  and so it has too few edges to contain any cycle. Therefore,  $G$  is a tree since it's also connected. Hence,  $m = n - 1$ . Since  $n \geq (g+2)/2$  by assumption, we have

$$g + 2 \leq 2n$$

i.e.,

$$gn - 2n - g + 2 \leq gn - 2g$$

i.e.,

$$(g-2)(n-1) \leq g(n-2)$$

i.e.,

$$m = n - 1 \leq \frac{g(n-2)}{g-2}$$

and the conclusion of the theorem holds.

Next, assume  $G$  has  $\geq g$  edges. Fix an embedding of  $G$  on the plane. For each region  $i$  (where  $1 \leq i \leq r$ ) of the plane graph  $G$ , let  $m_i$  be the number of edges on its boundary. Since  $G$  has at least  $g$  edges, has girth  $\geq g$ , and is connected, we see that  $m_i \geq g$  for each  $i$ . Thus  $\sum_{i=1}^r m_i \geq gr$ . Also,  $\sum_{i=1}^r m_i \leq 2m$  because, by Lemma A, each bridge contributes 1 to the sum and each nonbridge contributes 2 to the sum. Thus,  $gr \leq 2m$ ; hence,  $r \leq 2m/g$ . Combining this last inequality with Euler Identity we have

$$2 = n - m + r \leq n - m + \frac{2m}{g}$$

i.e.,

$$2g \leq gn - (g-2)m$$

i.e.,

$$(g-2)m \leq gn - 2g$$

i.e.,

$$m \leq \frac{g(n-2)}{g-2}$$

as desired. □

**Theorem** (Theorem 9.2 of CZ). *If  $G$  is a planar graph of order  $n$ , size  $m$ , and  $n \geq 3$ , then  $m \leq 3n - 6$ .*

*Proof.* Every graph has girth at least 3. Putting  $g = 3$  in the generalized Theorem 9.2 of CZ gives the result.  $\square$

**Theorem.** *If  $G$  is a bipartite planar graph of order  $n$ , size  $m$ , and  $n \geq 3$ , then  $m \leq 2n - 4$ .*

*Proof.* A bipartite graph has girth at least 4. Putting  $g = 4$  in the generalized Theorem 9.2 of CZ gives the result.  $\square$

**Theorem** (Corollary 9.3 of CZ). *Every planar graph contains a vertex of degree  $\leq 5$ .*

*Proof.* Let  $G$  be a planar graph of order  $n$  and size  $m$ . If  $n \leq 6$ , then every vertex has degree  $\leq 5$  and we are done. So assume  $n > 6$ . By Theorem 9.2,  $m \leq 3n - 6$ . Thus,

$$\frac{m}{n} \leq 3 - \frac{6}{n}$$

i.e.

$$\frac{2m}{n} \leq 6 - \frac{12}{n}$$

i.e.

$$\frac{2m}{n} < 6$$

since  $\frac{12}{n}$  is positive. The last inequality says that the average degree of  $G$  is  $< 6$ . Therefore, there exists at least a vertex whose degree does not exceed the average, i.e., some vertex  $v$  has  $\deg v \leq \frac{2m}{n} < 6$ , i.e.,  $\deg v \leq 5$ .  $\square$

**Theorem** (Corollary 9.4 of CZ).  *$K_5$  is nonplanar.*

*Proof.* By Theorem 9.2.  $\square$

**Theorem** (Theorem 9.5 of CZ).  *$K_{3,3}$  is nonplanar.*

*Proof.* By the fact that a bipartite planar graph satisfies  $m \leq 2n - 4$ .  $\square$

**Exercise.** Show that the Petersen graph is nonplanar by using the generalization of Theorem 9.2.

**Definition** A *subdivision*  $G'$  of a graph  $G$  is a graph that results from inserting one or more vertices of degree 2 into one or more edges of  $G$ .

**Theorem** (Kuratowski's Theorem). *Graph  $G$  is planar if and only if  $G$  contains no  $K_5$  or  $K_{3,3}$ , or subdivision of  $K_5$  or  $K_{3,3}$ , as a subgraph.*

**Exercise.** Show that the Petersen graph is nonplanar by using Kuratowski's Theorem.