## Bridge Lemma proof

## Lemma

If $\mathrm{e}=\mathrm{uv}$ is a bridge in a connected graph G, then G-e has exactly 2 connected components, one containing u and the other containing v .

## Proof(by contradiction)

1)In G-e, suppose $u$ and $v$ are in the same connected component $H$, since $H$ is connected, there must exist some path $p$ from $u$ to $v$. If $p$ exists in G-e, it must exist in G. With e and $p$ both connecting $u$ and $v$ in $G$, there exists a cycle containing e between $u$ and $v$. It contradicts the given condition that e is a bridge in G, so u and v must be in different components.

2 )we know G-e must contain at least 2 connected components. When connecting $u$ and $v$ by e, e, as an edge, can at most only join 2 vertices, which means e can only connect the 2 disconnected components that contain $u$ and v . If there are more than 2 disconnected components in G-e, the components without $u$ and $v$ will stay isolated. We know joining $u$ and $v$ by e in $\mathrm{G}-\mathrm{e}$ will give us G. In the above situation with more than 2 disconnected components in G-e, only components containing $u$ and $v$ will be connected by adding e while other components stay isolated, so G is disconnected, while the given condition says G is connected, so it contradicts. So only can there be exactly 2 disconnected components, one containing $u$ and the other containing $v$, could satisfy when joining $u$ and $v$ by e will give a connected graph.

Combining 1) and 2), lemma proved.

