

Study Questions for Greedy Algorithms

1. For the activity scheduling problem, give a set of classes, with their start and finish times, such that there are 2 optimum schedules—the first schedule can be constructed by the greedy algorithm, but the second schedule cannot.
2. Give a set of 4 classes such that the schedule obtained from earliest-finish-time-first algorithm differs from that obtained from latest-start-time-first algorithm.
3. Let S be an earliest-finish-time-first schedule for a collection of classes \mathcal{C} . Each class in \mathcal{C} starts and ends at some hour of the day, e.g. 3:00, not 3:20.
 - (i) The administration says that now every class of \mathcal{C} must end exactly 10 minutes earlier. No start times change. Is S still a valid schedule, and if so is it optimal?
 - (ii) The administration changes its mind. Each class must add an additional 10 minutes to its original ending time (for cleanup). No start times change. Is S still a valid schedule, and if so is it optimal?
4. Give one set of classes such that none of these greedy algorithms give an optimal schedule:
 - shortest-time first
 - longest-time first
5. Give pseudocode to implement the shortest-time-first algorithm in time $O(n^2)$. (Can you give one in time $O(n \log n)$?)
6. A set of items has weights 1, 2, 4, 5, 6.
 - (i) Draw a tree with minimum w.e.p.l. constructed by Huffman's algorithm.
 - (ii) Draw a tree for these items having minimum w.e.p.l. but cannot be constructed by Huffman's algorithm.
7. Professor Dull claims that If B is a binary tree such that swapping any of its two leaves increases the w.e.p.l. or keeps it the same, then B has minimum w.e.p.l. Give a counterexample to Dull's claim.
8. Given a set of items I with nonnegative weights, we want to find a binary tree having I as leaves and having **maximum** w.e.p.l.

- (i) Give a greedy algorithm for the problem.
- (ii) State a theorem that would prove this algorithm is correct.
- (iii) Give a swapping argument that proves the theorem.