Designing Finite Automata

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Feb 15, 2016

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- $1. \ \mbox{Given an FA},$ describe its language precisely.
- 2. Given a formal definition of a regular language, design an FA that recognizes it.

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2. Pinpoint the accept states.

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- 2. Pinpoint the accept states.
- 3. Draw the edges.

Assume $\Sigma=\{ \text{ a, b} \}.$ Design machines to recognize strings that

starts with a



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starts with a

ends with bb

Assume $\Sigma=\{ \text{ a, b }\}.$ Design machines to recognize strings that

- starts with a
- ends with bb
- starts with a and ends with b

Assume $\Sigma=\{ \text{ a, b} \}.$ Design machines to recognize strings that

- starts with a
- ends with bb
- starts with a and ends with b
- starts with a or ends with a

Assume $\Sigma = \{ a, b \}$. Design machines to recognize strings that

- starts with a
- ends with bb
- starts with a and ends with b
- starts with a or ends with a
- if it starts with a, then it doesn't end with a

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contains ab



- contains ab
- contains ab and ba



- contains ab
- contains ab and ba
- contains ab or ba



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- contains ab
- contains ab and ba
- contains ab or ba
- contains neither aa nor bb



contains (at least) an a



contains (at least) an a

contains exactly one a

- contains (at least) an a
- contains exactly one a
- contains (at least) two a's

- contains (at least) an a
- contains exactly one a
- contains (at least) two a's
- contains exactly two a's

- contains (at least) an a
- contains exactly one a
- contains (at least) two a's
- contains exactly two a's
- negate each of the above

Remarks

- We can always get the complement of a language of a DFA simply by exchanging the role of the accepting and the nonaccepting states. Does this technique work for NFAs?
- The Product Construction can be used to obtain the language of binary set operations (like union, intersection, set difference, symmetric difference, etc) of two regular languages.

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